

On Hedberg's Theorem: Proving and Painting

Nicolai Kraus

25/05/12

Reminder: Equality

Definitional Equality

“Real” decidable equality for typechecking, computation; e. g.
 $(\lambda a.b)x =_{\beta} b[x/a]$

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Propositional Equality

Equality needing a proof, i. e. a term of the identity type

Reminder: Identity Types

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Formation

$$\frac{a, b : A}{a \equiv b : \text{type}}$$

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Introduction

$$\frac{a : A}{\text{refl}_a : a \equiv a}$$

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Elimination (J)

$$\frac{P : (a, b : A) \rightarrow a \equiv b \rightarrow \text{Set} \\ m : \forall a. P(a, a, \text{refl}_a)}{J_{(a,b,q)} : P(a, b, q)}$$

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Computation (β)

$$J_{(a,a,\text{refl}_a)} =_{\beta} ma$$

Elimination: J versus K Eliminator J

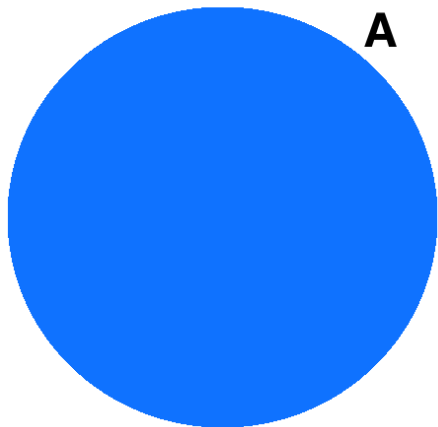
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 \end{array}$$

Eliminator K

$$\begin{array}{c}
 a : A \\
 P : a \equiv a \rightarrow \text{Set} \\
 P(\text{refl}_a) \\
 \hline
 K_q : P(q)
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Homotopic Model

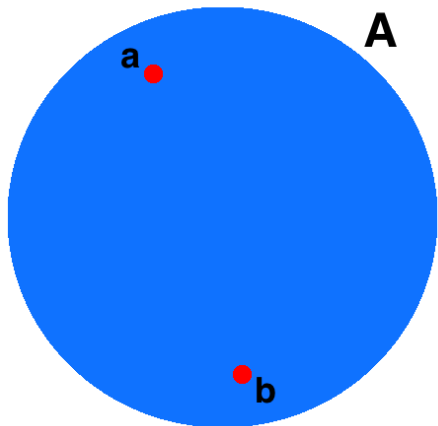
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Homotopic Model

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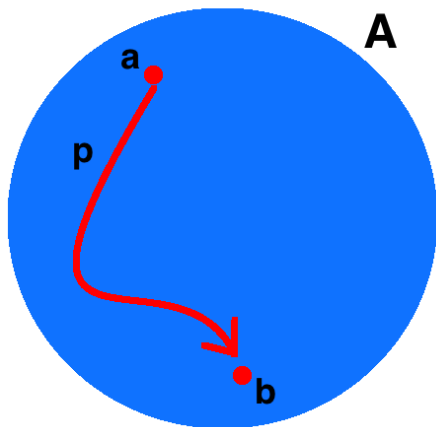


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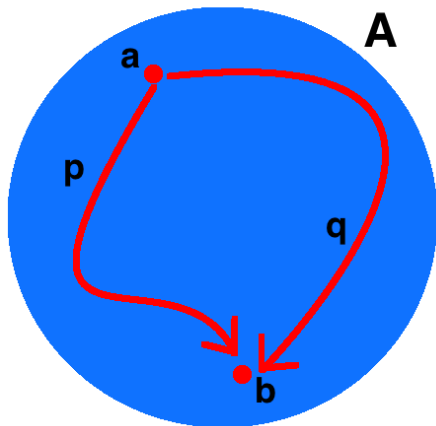


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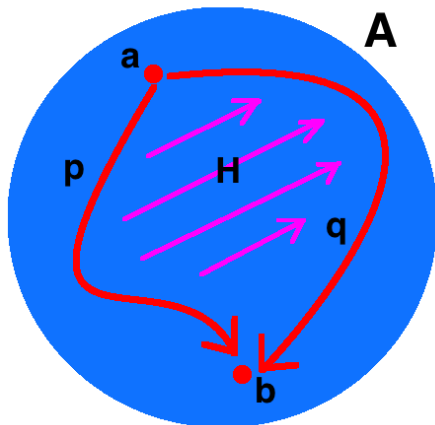
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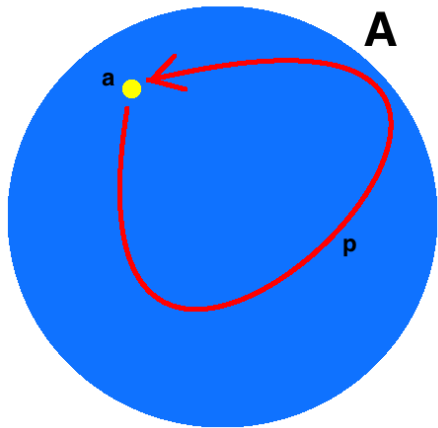
K revisited

Want to prove:

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Eliminator K

$$\frac{\begin{array}{l} a : A \\ P : a \equiv a \rightarrow \text{Set} \\ P(\text{refl}_a) \end{array}}{K_q : P(q)}$$



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only need $refl_a \equiv refl_a$

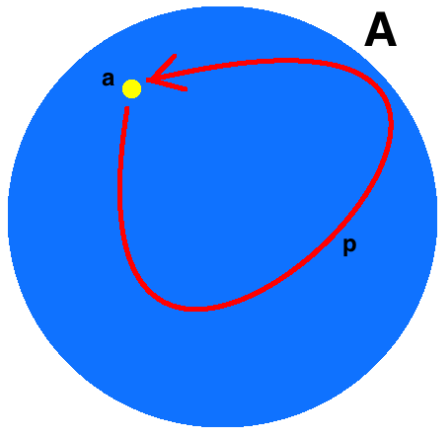
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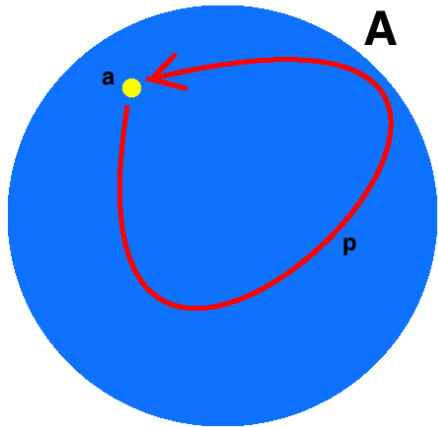
so, indeed $p \equiv refl_a$

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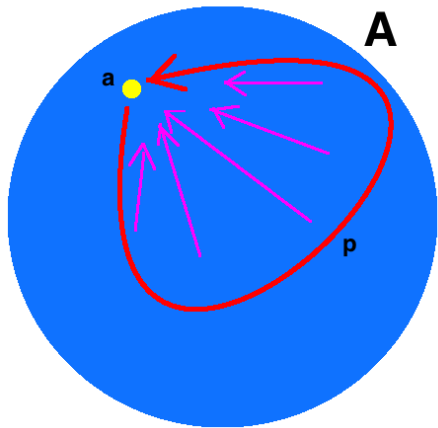
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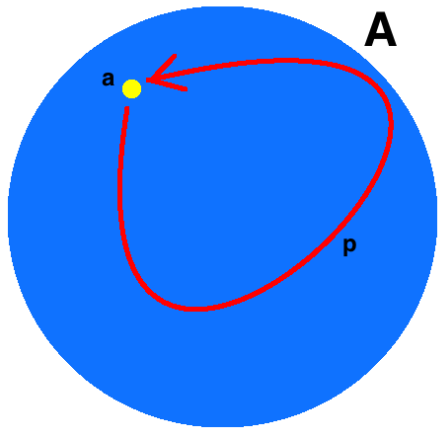
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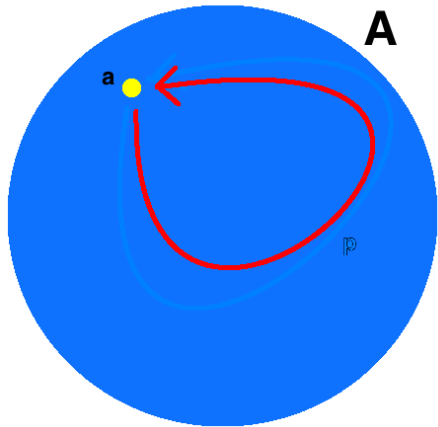
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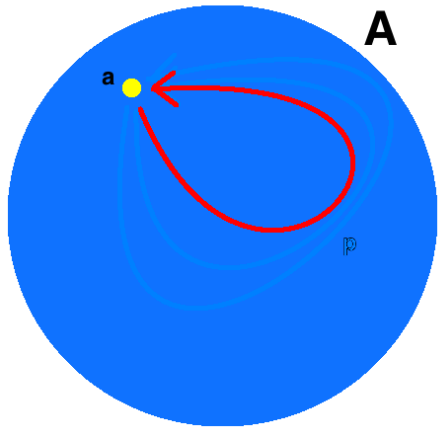
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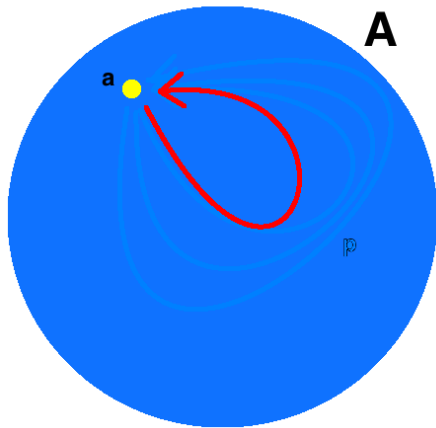
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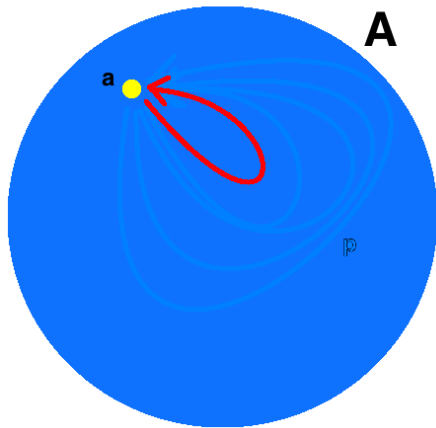
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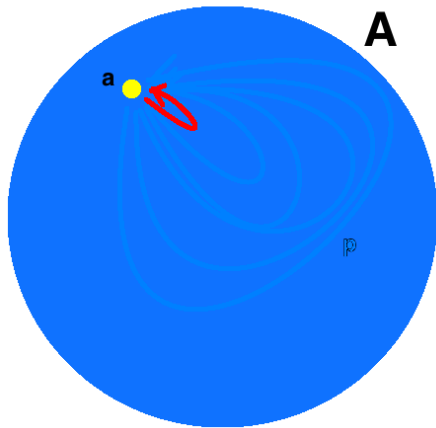
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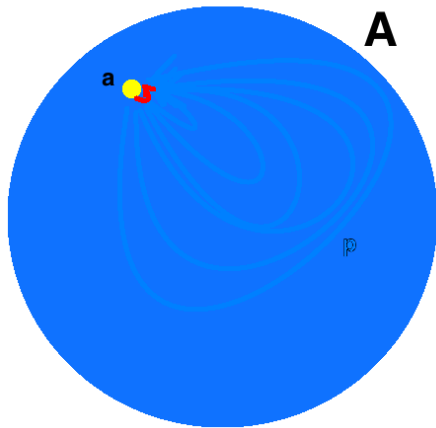
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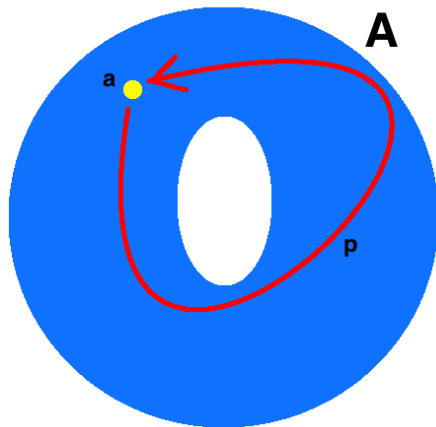
$$P : a \equiv a \rightarrow \text{Set}$$

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K revisited

Okay, but what now?



J revisited

Want: $(a, a, p) \equiv (a, a, refl_a)$

J shows:

$\forall(a, b, q) : \Sigma(a, b : A). a \equiv b,$

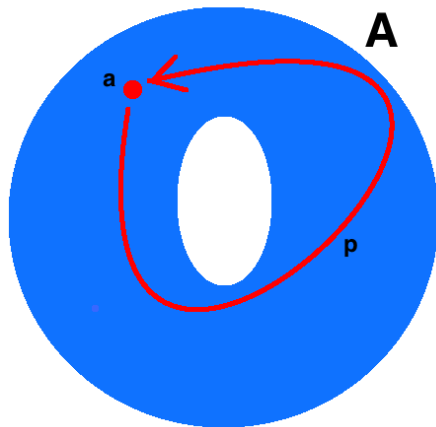
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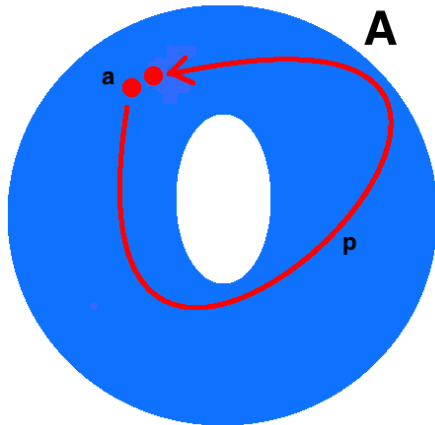
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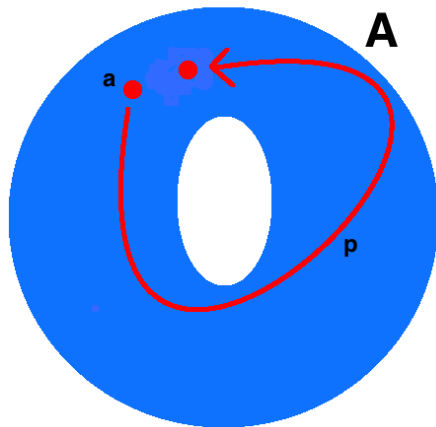
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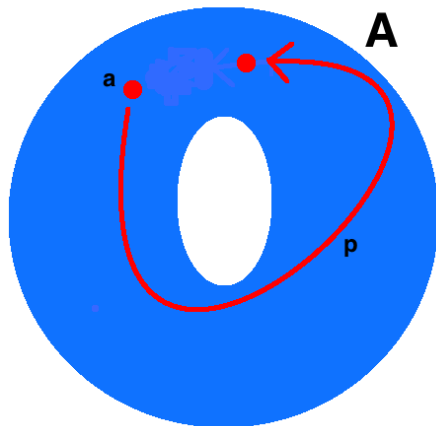
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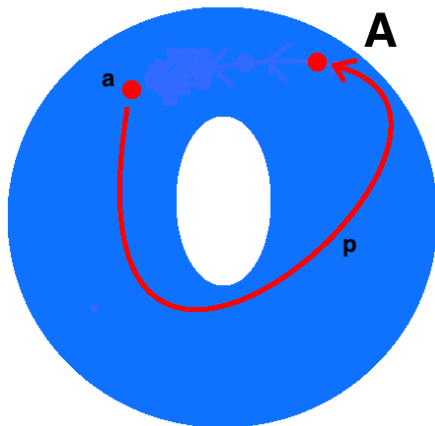
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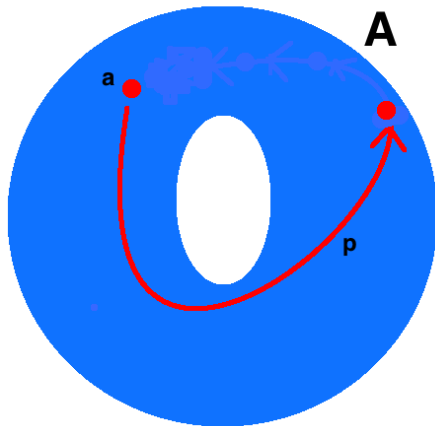
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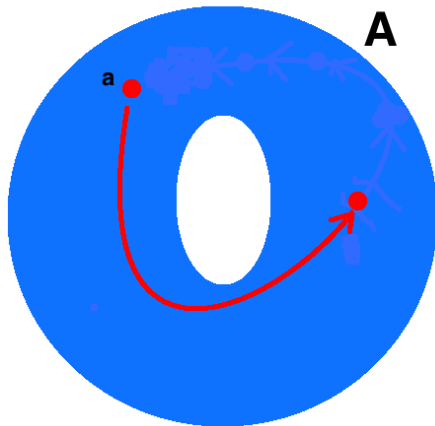
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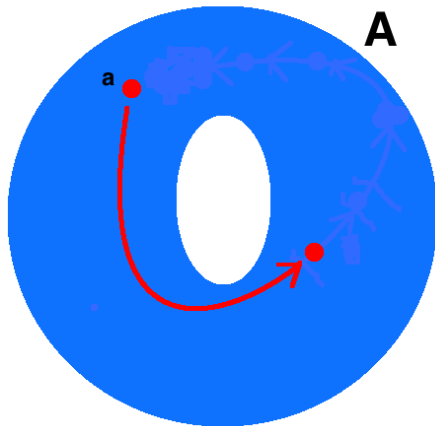
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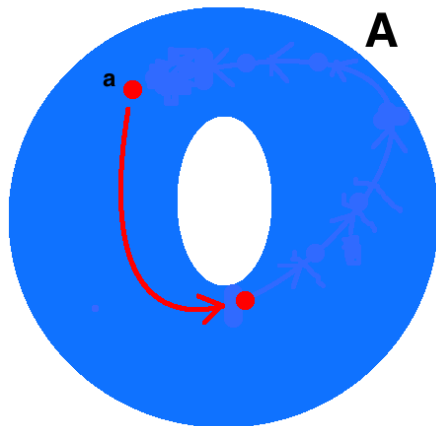
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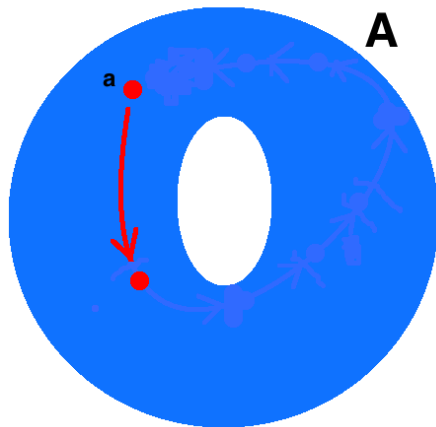
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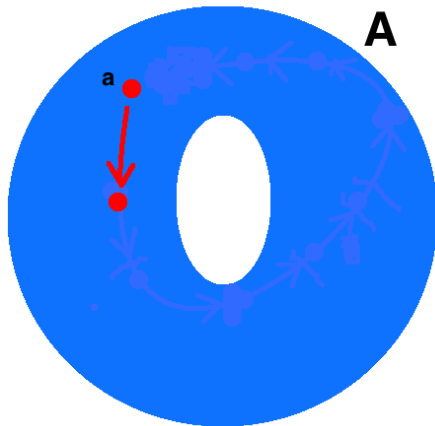
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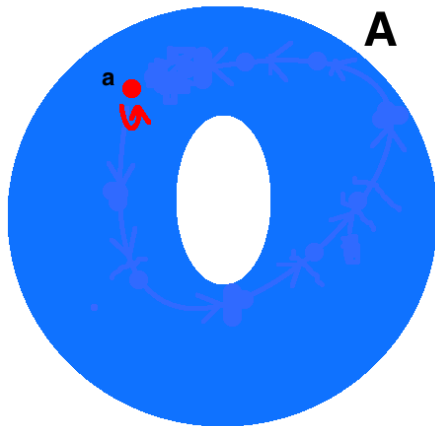
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J versus K revisited

Fix $a : A$.

J (variant of Paulin-Mohring)

To show $\forall (b, q) : \Sigma(b:A). a \equiv b \cdot P(b, q)$, just prove $P(a, refl_a)$.

K (reformulated)

$\forall q : a \equiv a \cdot q \equiv refl_a$

Hedberg's theorem

Fix a type A .

Decidable Equality

$$\text{DecidableEquality} := \forall a, b \ (a \equiv b + \neg a \equiv b)$$

Hedberg's theorem

$$\text{DecidableEquality} \longrightarrow K$$

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$DecidableEquality \rightarrow K$

Proof.

- Given $dec : (a, b : A) \rightarrow (a \equiv b + \neg a \equiv b)$.

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- Given $dec : (a, b : A) \rightarrow (a \equiv b + \neg a \equiv b)$.
- Given any $a, b : A$ and $p : a \equiv b$.

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$$\text{DecidableEquality} \longrightarrow K$$

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- Given $\text{dec} : (a, b : A) \rightarrow (a \equiv b + \neg a \equiv b)$.
- Given any $a, b : A$ and $p : a \equiv b$.
- $\text{dec } a \ b = \text{inl } q_1$ and $\text{dec } a \ a = \text{inl } q_2$ for some q_1, q_2 .

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- Claim: $p \equiv q_1 \circ q_2^{-1}$

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- Proof with J : Just show $\text{refl}_a \equiv q_2 \circ q_2^{-1}$. That's true!

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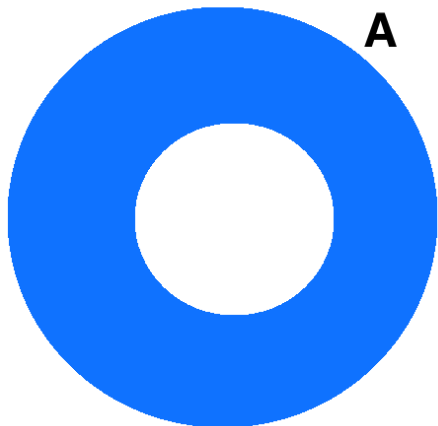
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- Claim: $p \equiv q_1 \circ q_2^{-1}$
- Proof with J : Just show $\text{refl}_a \equiv q_2 \circ q_2^{-1}$. That's true!
- Special case: If $p : a \equiv a$, then $p \equiv \text{refl}_a$. \square

Corollary¹: The Circle type does not have decidable equality

$dec : (a, b : A) \rightarrow$
 $(a \equiv b + \neg a \equiv b)$



¹If you believe me that K does not hold for the Circle type

Nearly uncountable many things to be done . . .

- Higher Inductive Types (see Mike Shulman's work)
- Model construction with modern abstract (not point-set) homotopy theory
- Constructive Simplicial Sets (the combinatorial version of what I have shown; see Thierry Coquand's / Simon Huber's work)
- Univalent foundations / Univalence ("alternative" to K) in general (see Voevodsky)
- . . . and possible computational properties (Thorsten?)

THANK YOU!